

TRANSACTION ASSISTING APPARATUS AND RECORDING MEDIUM

BACKGROUND OF THE INVENTION

(1) Field of the Invention

5 The present invention relates to a transaction assisting apparatus and a recording medium, and more particularly, to a transaction assisting apparatus for assisting transactions wherein sell/buy orders from a plurality of customers are displayed at display units and when a deal for a
10 certain order is offered, an agent places an order for corresponding goods to thereby conduct the deal of goods, and to a computer-readable recording medium storing a program for causing a computer to perform such a transaction assisting process.

15 (2) Description of the Related Art

 In recent years, a variety of transaction assisting apparatus are used with a view to ensuring reliability and swiftness of transactions in goods etc.

 FIG. 22 is a block diagram exemplifying a configuration
20 of a transaction assisting system comprising a transaction assisting apparatus for assisting transactions in bonds, terminal units, etc.

 In the figure, a transaction assisting apparatus 1 accepts and stores sell and buy orders etc. for bonds, which are
25 input from a terminal unit 2 owned by an agent, performs a predetermined process on the orders, and makes the orders displayed at the terminal unit 2 as well as at display units 3-1

to 3-n installed at customers' side.

A telephone 4 installed at the agent's side is connected via a leased line (hot line) 6 to customers' telephones 5-1 to 5-n.

5 Let it be assumed that, in this transaction assisting system, a customer A, for example, has requested, using the telephone 5-1, a sell order for bonds with the circular issue (a sort of serial number) "183" to the agent having the telephone 4.

10 On receiving the request, the agent operates the terminal unit 2 to input the sell order for "183". The transaction assisting apparatus 1 detects the input of the sell order for the bonds "183", stores the input sell order in a storage device therein, and displays the sell order at the agent's
15 terminal unit 2 as well as at the display units 3-1 to 3-n.

FIG. 23 shows an example of a screen displayed in this case at the terminal unit 2 and the display units 3-1 to 3-n. In the illustrated example, sell and buy orders for the bonds "183" are shown in list form. Also, in the screen shown in the
20 figure, sell orders and buy orders are shown in left- and right-hand regions of the list, respectively. "ORDER NO." is a number assigned so as to specify each sell or buy order, and "RATE" indicates the yield of the bond concerned. "AMOUNT" indicates an amount ordered, and "1 BILLION" means one billion
25 yen. The display screen as shown in the figure is called board screen.

In the illustrated example, only one order is shown,

but in cases where there are a plurality of orders, sell orders and buy orders are respectively shown in order of rate such that the order with the highest rate (hereinafter referred to as best rate) is listed at top.

5 If, looking at the display screen, a customer B, for example, puts in a buy order for "183" with the agent by using the telephone 5-2, the agent enters the requested buy order through the terminal unit 2. Consequently, the transaction assisting apparatus 1 performs a contracting process (process
10 executed when a deal is to be closed) with respect to the sell order for "183".

 As a result, the deal for "183" is conducted, whereupon the transaction assisting apparatus 1 deletes the sell order for "183" from the storage device and also deletes the display item
15 corresponding to the sell order for "183" from the display screens of the terminal unit 2 and the display units 3-1 to 3-n. If there are other sell orders remaining, display items corresponding thereto are appropriately rearranged according to rates.

 The use of the aforementioned transaction assisting
20 apparatus 1 permits reliable and speedy transactions.

 In the dealing business, a so-called chain deal is conducted in which one purchases somebody else's goods on condition that he or she sells his/her own goods.

 In the case of bonds, for example, chain deal is
25 conducted in such a manner that one purchases somebody else's "200" bonds if a buy order is offered for his/her "183" bonds, or conversely, one sells his/her "183" bonds if a sell order for

"200" bonds is offered by somebody else.

Let us consider the case where such chain deal is conducted with the use of the transaction assisting system shown in FIG. 22. If a customer puts in the chain order as mentioned
5 above, for example, the agent enters a sell order for "183" through the terminal unit 2, whereupon the sell order for "183" is displayed at the terminal unit 2 and the display units 3-1 to 3-n.

This chain order is a sell order premised on the
10 presence of a chained order (in this example, "200") on which the chain order chains, and accordingly, the agent must always make sure whether there is a sell order for "200", which is a chained order, and must immediately suspend the sell order for "183" when there is no chained order present.

Thus, when such a chain order is accepted, it is
15 necessary for the agent to always direct his or her attention to chained orders, making the work complicated.

Also, when a buy order is offered for the chain order "183", a buy order for the chained order "200" must be placed
20 immediately in a manner interlocked with such a buy order, which increases the burden imposed on the agent.

Further, in cases where chain orders are chained (for example, the above order for "200" also is a chain order which chains on "201"), the aforementioned operation needs to be
25 repeatedly performed, which often makes it difficult to carry out the process reliably and speedily.

In the transactions in used cars, real estate, etc.,

so-called replacement deal is a customary practice wherein one purchases somebody else's goods or real estate on condition that a buy order is offered for his or her goods or real estate, or conversely, one sells his/her goods or real estate on condition that a sell order for goods or real estate is offered by somebody else. Problems similar to those described above also arise in this kind of transactions.

Further, in the case of transactions in used cars, real estate, etc., a person who has to pay taxes changes from one to another at the time when the ownership and thus the title are transferred. Accordingly, if the previous owner has paid taxes extra, it is necessary that the overpayment be refunded and also that the deficit be collected from the current owner, thus requiring complicated procedures.

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SUMMARY OF THE INVENTION

An object of the present invention is to provide a transaction assisting apparatus which makes it possible to reliably and speedily conduct chain orders and replacement orders in transactions in goods etc.

Another object of the present invention is to provide a transaction assisting apparatus which makes it possible to reliably and easily carry out tax-related procedures etc. that accompany the transfer of ownership of goods etc.

To achieve the above objects, there is provided a transaction assisting apparatus for assisting transactions in goods. The transaction assisting apparatus comprises chain

order input means permitting input of a chain order, the chain order being an order requesting, on condition of sale or purchase of first goods, purchase or sale of second goods different from the first goods, order detecting means for detecting an order
5 for purchase or sale of the first goods with respect to which the chain order is placed, and chain order processing means, responsive to the detection of a buy or sell order with respect to a predetermined chain order by the order detecting means, for performing a selling or purchasing process with respect to the
10 first goods and performing a purchasing or selling process with respect to the second goods.

The above and other objects, features and advantages of the present invention will become apparent from the following description when taken in conjunction with the accompanying
15 drawings which illustrate preferred embodiments of the present invention by way of example.

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 is a diagram illustrating the principle of
20 operation according to the present invention;

FIG. 2 is a diagram showing an example of information stored in chain order storing means appearing in FIG. 1;

FIG. 3 is a diagram showing an example of a display screen displayed at a display unit appearing in FIG. 1;

25 FIG. 4 is a diagram showing an example of a display screen which is displayed when a sell order for a chained issue is newly offered in FIG. 3;

FIG. 5 is a block diagram exemplifying a configuration according to one embodiment of the present invention;

FIG. 6 is a flowchart showing an example of a process executed in the embodiment shown in FIG. 5;

5 FIG. 7 is a flowchart illustrating details of a "DISPLAY PROCESS" shown in FIG. 6;

FIG. 8 is a flowchart illustrating details of a "CANCELLATION PROCESS" shown in FIG. 6;

10 FIG. 9 is a flowchart illustrating details of a "RATE FOLLOW-UP PROCESS" shown in FIG. 6;

FIG. 10 is a flowchart illustrating details of a "CONTRACTING PROCESS" shown in FIG. 6;

FIG. 11 is a block diagram exemplifying a configuration according to a second embodiment of the present invention;

15 FIG. 12 is a diagram illustrating a vehicle replacement deal;

FIG. 13 is a flowchart showing an example of a process executed when an order is input in the embodiment shown in FIG. 11;

20 FIG. 14 is a diagram showing an example of an order input from each terminal unit in the embodiment shown in FIG. 11;

FIG. 15 is a diagram showing another example of an order input from each terminal unit in the embodiment shown in FIG. 11;

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FIG. 16 is a diagram showing an example of a screen displayed at each terminal unit in the embodiment shown in FIG.

11;

FIG. 17 is a flowchart showing an example of a process executed when an input order is processed in the embodiment shown in FIG. 11;

5 FIG. 18 is flowchart illustrating in detail a "BUY ORDER PREFERENCE PROCESS" shown in FIG. 17;

FIG. 19 is flowchart illustrating in detail a "SELL ORDER PREFERENCE PROCESS" shown in FIG. 17;

10 FIG. 20 is a diagram showing another example of a screen displayed at each terminal unit in the embodiment shown in FIG. 11;

FIG. 21 is a flowchart illustrating an example of an "ADJUSTMENT PROCESS" shown in FIG. 17;

15 FIG. 22 is a diagram exemplifying a configuration of a conventional transaction assisting system; and

FIG. 23 is a diagram showing an example of a display screen displayed at a terminal unit and display units of the system shown in FIG. 22.

20 DESCRIPTION OF THE PREFERRED EMBODIMENTS

Embodiments of the present invention will be hereinafter described with reference to the drawings. It is to be noted that in the specification, "goods" represent "property which can be dealt in" and include not only "movable property" but "fixed property" and "intangible property" and that in the
25 claims, "chain order" includes "replacement order."

FIG. 1 illustrates the principle of operation of a

transaction assisting apparatus according to the present invention.

In the figure, chain order input means 50a permits input of a chain order, which is an order requesting, on condition
5 of sale or purchase of first goods, purchase or sale of second goods different from the first goods.

Chain order storing means 50b stores the chain order input from the chain order input means 50a. The chain order storing means 50b stores, in addition to chain orders, ordinary
10 sell orders and buy orders.

Order detecting means 50c detects a buy or sell order for the first goods which are a target of a predetermined chain order stored in the chain order storing means 50b.

When such a buy or sell order with respect to the
15 predetermined chain order is detected by the order detecting means 50c, chain order processing means 50d performs a selling or purchasing process with respect to the first goods and a purchasing or selling process with respect to the second goods.

Deleting means 50e deletes, from the chain order
20 storing means 50b, a chain order with respect to which processing has been completed by the chain order processing means 50d.

If the price of the second goods has changed because of a new sell or buy order etc., price changing means 50f changes the price of the first goods in a manner interlocked with the
25 change of the price of the second goods. When it is unnecessary to change prices, the price changing means 50f supplies the chain orders stored in the chain order storing means 50b directly to

display means 50g.

The display means 50g causes a display unit 60 to display the chain order of which the price has been appropriately changed by the price changing means 50f.

5 The operation in accordance with the aforementioned principle will be now described. The following description is given taking as an example a chain deal in which one purchases somebody else's goods on condition of sale of his or her goods.

Assuming that a customer A places a chain order for
10 purchase of bonds "200" on condition of sale of his/her bonds "183", the chain order is entered by, for example, an agent, through the chain order input means 50a. In the following, the bonds (in this example, "183") which are to be sold by the chain order are called chain issue, and the bonds (in the example,
15 "200") to be purchased are called chained issue.

The chain order input from the chain order input means 50a is stored in the chain order storing means 50b.

FIG. 2 shows, by way of example, a chain order stored in the chain order storing means 50b. In the illustrated example
20 is shown information indicative of a chain order specifying the aforementioned "200" as the chained issue (on which the chain issue chains) and "183" as the chain issue. "AMOUNT" indicates an amount of the chain issue to be sold and, in this example, "1 BILLION" yen.

25 "RATE" is a value determining the selling rate of "183", and the illustrated example indicates that "0.1" is to be added to the rate of the chained issue "200". Accordingly, if the rate

of "200" is "1.05", the rate of "183" is "1.15".

"DESIGNATION RANGE" indicates that if the best rate of the chained issue is within this range, the corresponding chain order is valid. In the illustrated example, if a buy order for the chain issue "183" is offered and if the best rate of the chained issue "200" is within the range from "1.00" to "1.20", a buy order is placed for the chained issue "200".

The price changing means 50f reads out the chain order stored in the chain order storing means 50b, and if the price of the chained issue has changed, it changes the price of the chain issue and supplies the changed price to the display means 50g. If there has been no change in price, the information is supplied to the display means 50g without being modified.

The display means 50g causes the display unit 60 to display the information supplied from the price changing means 50f. In this case, the display means 50g assigns an order number to each order, and also affixes information indicative of chain order (as described in detail later) to an order if the order is a chain order.

FIG. 3 shows an example of a display screen displayed at the display unit 60 as a result of the aforementioned process. In the illustrated example is shown information relating to "183" and "200". With regard to "183", the information corresponding to the previously input chain order is shown with the order number "8" assigned thereto. "H" appearing in the lower right corner of the cell with the order number "8" indicates that the order concerned is a chain order, and is information affixed by the

display means 50g. In "RATE" is shown a value (= 1.150) which is obtained by adding "0.1" to the rate of the sell order with the order number "2" displayed as the information relating to "200" below the information on "183".

5 Basically, chain order is a sell/buy order for chained goods, and thus the sell/buy order for other goods should also be displayed on the board screen. If, however, such a sell/buy order for other goods is displayed, a sell order, for example, can possibly be placed for the buy order, and therefore, such
10 a sell/buy order is shown temporarily after the deal of the chain issue is transacted, and not at the time when the order is placed.

 Let it be assumed that a customer B, for example, puts in a buy order for "183" while viewing the display screen as shown in FIG. 3, the order detecting means 50c makes reference to the
15 chain order storing means 50b to determine whether there is a chain order which chains on the input buy order. If, as a result, it is found that there is an applicable chain order, the order detecting means 50c acquires the corresponding information from the chain order storing means 50b and supplies the acquired
20 information to the chain order processing means 50d.

 While referring to the chain order supplied from the order detecting means 50c, the chain order processing means 50d performs a selling process with respect to the chain issue and a purchasing process with respect to the chained issue.

25 In this example, "183" is the chain issue and "200" is the chained issue; accordingly, the selling process is performed with respect to "183" and then the purchasing process

is performed with respect to "200". As a result, "183" is sold to the customer B and "200" is sold to the customer A who placed the chain order, whereby a contract is made.

After the processing of the chain order is completed,
5 the chain order processing means 50d notifies the deleting means 50e of completion of the processing. On receiving such notification from the chain order processing means 50d, the deleting means 50e deletes the corresponding chain order from the chain order storing means 50b.

10 Consequently, the chain order shown in FIG. 2 is deleted from the chain order storing means 50b and the order with the order number "8", shown in FIG. 3, disappears from the screen of the display unit 60.

At this time, the display means 50g causes the display
15 unit 60 to temporarily display a message informing that a buy order for the chain issue has been accepted.

Such a buy order for the chain issue may alternatively be displayed on the board screen while specifying that the order concerned is a buy order related to the chain order.

20 As described above, the transaction assisting apparatus 50 according to the present invention automatically conducts chain orders, whereby the burden imposed on the agent etc. can be mitigated.

In the foregoing description, fluctuation in price of
25 the chained issue and other factors are not taken account of; however, in actual transactions, price is subject to fluctuation and also a sell order for the chained issue may be canceled. The

following explains a process executed in such situations.

Let it be assumed that a sell order is newly put in for the chained issue "200", for example, and that the rate of this sell order is the highest (best rate) among those then available. Chain order should chain on the chained issue with the best rate, and therefore, if the best rate has changed, the rate of the chain issue needs to be varied so as to follow up such change.

Assuming that a new sell order with the rate "1.055" is placed for the chained issue "200", for example, this order is displayed at the top of the display area associated with "200", as shown in FIG. 4, since it has the best rate.

In this case, the price changing means 50f detects the change in the best rate of the chained issue, changes the rate of the chain issue "183", and supplies the changed rate to the display means 50g. Specifically, in this example, since the rate of the chain issue is the sum of the rate of the chained issue and "0.1", as shown in FIG. 2, "1.155" is calculated as the new rate and is supplied to the display means 50g.

The display means 50g generates a new order number "12" for the chain order whose rate has been changed, and updates the contents of information displayed at the display unit 60.

As a result, a chain order having the order number "12" and the rate "1.155" is displayed at the display unit 60, as shown in FIG. 4.

As described above, with the transaction assisting apparatus 50 according to the present invention, the rate of the

chain issue is varied automatically in a manner interlocked with fluctuation in the rate of the chained issue, whereby the burden imposed on the agent etc. can be mitigated.

5 The following describes a process executed in cases where the best rate of the chained issue becomes outside the designation range, the chained issue is sold out, or the sell order has been canceled and there no longer exists a chained issue to be dealt in.

10 If the best rate of the chained issue changes to "1.21", for example, then it is outside the designation range because the chain order for the chain issue "183" has the designation range from "1.00" to "1.20", as shown in FIG. 2.

15 In such a case, the display means 50g immediately stops the output of this chain order, so that the corresponding information disappears from the screen of the display unit 60.

Similarly, in cases where a sell order for the chained issue has gone, the display of the chain order is stopped by the display means 50g, and thus the corresponding information disappears from the screen.

20 The display of the chain order is stopped only temporarily, and if the best rate of the chained issue comes to fall within the designation range or a new sell order is placed for the chained issue, the chain order is again displayed on screen.

25 As described above, with the transaction assisting apparatus 50 according to the present invention, in cases where a sell order for the chained issue has gone or the best rate of

the chained issue becomes outside the designation range, the display of the chain order is temporarily stopped. It is therefore unnecessary for the agent to stop or restart the display of the chain order while watching the conditions for display, 5 whereby the burden imposed on the agent can be mitigated.

In the case where an order for the chained issue also is a chain order, the chain order processing means 50d performs a similar process also with respect to the chained issue, so that a chain of chain orders can be coped with. If, in the example 10 stated above, an order for the issue "200", on which "183" chains, is a chain order specifying "201" as the chained issue, for example, the purchasing process for "201" is additionally executed following the aforementioned process.

It is therefore possible to further mitigate the burden 15 imposed on the agent and to carry out transactions reliably and speedily.

Referring now to FIG. 5, an exemplary configuration according to one embodiment of the present invention will be described.

20 FIG. 5 is a block diagram exemplifying a configuration according to the embodiment of the present invention.

In the figure, a transaction assisting apparatus 1 accepts and stores sell orders, buy orders, etc. input through a terminal unit 2 owned by an agent, subjects the input orders 25 etc. to a predetermined process such as sorting, and causes the terminal unit 2 and display units 3-1 to 3-n installed at customers' side to display the processed data.

As shown in the figure, the transaction assisting apparatus 1 comprises a CPU (Central Processing Unit) 1a, a ROM (Read Only Memory) 1b, a RAM (Random Access Memory) 1c, an HDD (Hard Disk Drive) 1d, I/Fs (Interfaces) 1e and 1f, and a bus 1g.

5 The CPU 1a controls various sections of the apparatus in accordance with programs etc. stored in the ROM 1b, the RAM 1c and the HDD 1d, and performs various operations.

 The ROM 1b stores basic programs executed by the CPU 1a, as well as data, etc.

10 The RAM 1c temporarily stores programs to be executed by the CPU 1a, data derived in the middle of calculations, etc.

 The HDD 1d stores various programs executed by the CPU 1a, and also stores sell orders, buy orders, chain orders, etc. input through the terminal unit 2.

15 The I/F 1e is an interface for the terminal unit 2, and converts data format from one to another to permit exchange of information between the terminal unit 2 and the transaction assisting apparatus 1.

 The I/F 1f supplies image information simultaneously
20 to the display units 3-1 to 3-n and causes the display units to display a board screen as shown in FIG. 23. The contents of display of the terminal unit 2 are the same as those of the display units 3-1 to 3-n, and thus the agent and the customers share identical information.

25 The bus 1g interconnects the CPU 1a, the ROM 1b, the RAM 1c, the HDD 1d, and the I/Fs 1e and 1f, to permit information to be exchanged with one another.

A telephone 4 installed at the agent's side is connected via a leased line (hot line) 6 to telephones 5-1 to 5-n installed at the customers' side.

The operation of the above embodiment will be now described with reference to FIG. 6.

FIG. 6 is a flowchart illustrating an example of a process executed when a chain order is placed in the embodiment shown in FIG. 5. When the process is started, the following steps are executed.

[S1] The CPU 1a determines whether or not a chain order has been put in. If a chain order has been put in, the flow proceeds to Step S2; otherwise the flow proceeds to Step S3.

Namely, when a chain order has been input through manipulation of the terminal unit 2, the CPU 1a executes Step S2; otherwise it executes Step S3.

[S2] The CPU 1a records the chain order input in Step S1 in a predetermined file (reserved order file; see FIG. 2) stored in the HDD 1d.

[S3] The CPU 1a performs, with respect to the terminal unit 2 and the display units 3-1 to 3-n, a "DISPLAY PROCESS" whereby the chain orders stored in the reserved order file shown in FIG. 2 are suitably displayed at these units.

Details of the process will be described later with reference to FIG. 7.

[S4] If a certain chain order fails to fulfill required conditions, the CPU 1a executes a "CANCELLATION PROCESS" whereby the chain order concerned is canceled.

Details of this process will be described later with reference to FIG. 8.

[S5] If the best rate of the chained issue has changed, the CPU 1a executes a "RATE FOLLOW-UP PROCESS" to cause the rate
5 of the chain issue to follow up the best rate in a manner interlocked therewith.

The process will be described in detail later with reference to FIG. 9.

[S6] If a buy order is put in for the chain issue with
10 respect to which a chain order has been placed, the CPU 1a executes a "CONTRACTING PROCESS", which is a series of operations required to transact the chain order.

Details of this process will be described later with reference to FIG. 10.

[S7] The CPU 1a determines whether or not the process
15 is to be ended. If the process should not be ended, the flow returns to Step S1, otherwise the process is ended (END).

For example, if execution of a process other than the illustrated process is requested, the CPU 1a terminates the
20 process; otherwise the flow returns to Step S1.

Referring now to FIG. 7, the "DISPLAY PROCESS" shown in FIG. 6 will be described in detail. Upon start of the process shown in the flowchart, the following steps are executed.

[S10] The CPU 1a determines whether or not a chain order
25 is recorded in the reserved order file stored in the HDD 1d. If a chain order is recorded, the flow proceeds to Step S11; otherwise the flow returns to the original process (RETURN).

[S11] The CPU 1a acquires one predetermined chain order from the reserved order file.

[S12] The CPU 1a determines whether or not there remains the amount of the chained issue requested by the chain order acquired in Step S11, that is, whether or not there is a sufficient amount of the chained issue with the best rate. If there is a sufficient amount of the chained issue, the flow proceeds to Step S13; if not, the flow proceeds to Step S17.

[S13] The CPU 1a determines whether or not the rate of the best-rate chained issue falls within the designation range. If the rate is within the designation range, the flow proceeds to Step S14, and if not, the flow proceeds to Step S17.

[S14] The CPU 1a determines whether or not the chain order being processed is already displayed. If the chain order is already displayed, the flow proceeds to Step S17; if not, the flow proceeds to Step S15.

[S15] The CPU 1a issues an order number.

[S16] The CPU 1a causes the terminal unit 2 and the display units 3-1 to 3-n to display the chain order being processed on their board screen.

Usually, since chain order involves a sell/buy order (in this example, buy order) for chained goods, the sell/buy order is displayed also on the board screen associated with the chained goods. However, such a buy order is not displayed because it is not transacted singly.

[S17] The CPU 1a refers to the reserved order file to determine whether or not there is an unprocessed chain order.

If an unprocessed chain order exists, the flow returns to Step S11 to repeat the same process as described above; if not, the flow returns to the original process (RETURN).

According to the above process, if a chain order is
5 recorded in the reserved order file and if it fulfills the chain order display conditions (if there remains a sufficient amount of the chained issue, if the rate is within the designation range, and if the chain order is not displayed yet), the chain order is displayed on the board screens with an order number assigned
10 thereto. If the chain order thereafter fails to satisfy the display conditions, it is deleted from the board screens in the "CANCELLATION PROCESS" or the "RATE FOLLOW-UP PROCESS" described below.

Referring now to FIG. 8, the "CANCELLATION PROCESS"
15 shown in FIG. 6 will be described in detail. When the process shown in the flowchart is started, the following steps are executed.

[S20] The CPU 1a determines whether or not a sell order has been canceled. If a sell order has been canceled, the flow
20 proceeds to Step S21; otherwise the flow returns to the original process (RETURN).

[S21] The CPU 1a refers to the reserved order file to determine whether or not there is a chain order which chains on the canceled sell order. If such a chain order exists, the flow
25 proceeds to Step S22, and if not, the flow returns to the original process (RETURN).

[S22] The CPU 1a deletes the applicable chain order

from the board screens, and then resumes the original process (RETURN).

According to the above process, if a sell order is canceled and if a certain chain order chains on the sell order, the chain order is deleted from the board screens. Needless to say, the canceled sell order is also deleted from the board screens by another process, not shown.

Referring now to FIG. 9, the "RATE FOLLOW-UP PROCESS" shown in FIG. 6 will be described in detail. Upon start of the process shown in the flowchart, the following steps are executed.

[S30] The CPU 1a determines whether or not a new sell order has been placed. If a sell order has been newly placed, the flow proceeds to Step S31; otherwise the flow returns to the original process (RETURN).

[S31] The CPU 1a determines whether or not the sell order is an order with the best rate (best-rate order). If the sell order is a best-rate order, the flow proceeds to Step S32; if not, the flow returns to the original process (RETURN).

[S32] The CPU 1a refers to the reserved order file to determine whether or not there exists a chain order which chains on the best-rate order. If such a chain order exists, the flow proceeds to Step S33; otherwise the flow returns to the original process (RETURN).

[S33] The CPU 1a calculates a new rate corresponding to the best rate. Specifically, the CPU 1a adds the rate specified when the chain order was placed (the rate shown in FIG. 2) to the rate of the chained issue whose rate has been judged

to be highest (best rate) in Step S31, thereby calculating the rate of the chain issue.

[S34] The CPU 1a deletes the chain order in question from the board screens.

5 [S35] The CPU 1a looks up the designation range stored in the reserved order file, and determines whether or not the new rate falls within the designation range. If the new rate is within the designation range, the flow proceeds to Step S36; if not, the flow returns to the original process (RETURN).

10 [S36] The CPU 1a issues a new order number.

[S37] The CPU 1a displays, on the board screens, the chain order to which the new order number and the new rate have been assigned.

15 According to the above process, when the best rate of the chained issue has changed, the corresponding chain issue is deleted from the board screens, and the chain order is again displayed with a new order number and a new rate assigned thereto. Also, if the new rate does not fall within the designation range, the corresponding chain order is deleted from the board screens.

20 Lastly, referring to FIG. 10, details of the "CONTRACTING PROCESS" shown in FIG. 6 will be described. When the process shown in the flowchart is started, the following steps are executed.

[S40] The CPU 1a determines whether or not a buy order
25 has been placed. If a buy order has been placed, the flow proceeds to Step S41; otherwise the flow returns to the original process (RETURN).

[S41] The CPU 1a refers to the reserved order file, and if the buy order placed is a buy order for a chain order, the flow proceeds to Step S42; otherwise the flow proceeds to Step S48.

5 [S42] The CPU 1a automatically conducts a deal with respect to the chain order concerned. Specifically, a selling process for the chain issue and a purchasing process for the chained issue are performed.

[S43] The CPU 1a deletes the corresponding chain order
10 from the board screens.

[S44] The CPU 1a deletes the corresponding chain order from the reserved order file.

[S45] The CPU 1a deletes the corresponding chained order from the board screens. At this time, the CPU 1a
15 temporarily displays, on the board screens, information indicating that a buy order for the chained goods has been conducted.

[S46] The CPU 1a deletes the corresponding chained order from the reserved order file.

20 In the case where the amount of the chained order is greater than that of the chain order, the remainder of the chained order is displayed with a new order number issued (for example, the number "1-2" is issued to indicate that the order is derived from the original order with the order number "1").

25 [S47] The CPU 1a determines whether or not the chained order is a chain order. If the chained order is a chain order, the flow returns to Step S42 to repeat the same process as

described above; if not, the flow returns to the original process (RETURN).

Namely, in the case where chain orders are chained, the flow returns to Step S42 and the aforementioned process is repeated, so that a chain of chain orders can be coped with.

In order to avoid the process becoming too complicated, the process may be interrupted in the event more than a predetermined number of chain orders are chained.

[S48] The CPU 1a conducts an ordinary deal. Namely, since the buy order placed is not for a chain order and is an ordinary order (outright order), an ordinary process is carried out with respect to the outright order and the original process is resumed (RETURN).

According to the above process, if a buy order is placed and if the order placed is for a chain order (chain issue), a selling process for the chain issue and a purchasing process for the chained issue are automatically carried out, and the results of the deals conducted are reflected in the board screens.

As described above, according to the embodiment of the present invention, chain orders are automatically transacted, permitting reliable and speedy transactions.

In the event the best rate of the chained issue has changed, the rate of the chain issue is automatically varied so as to follow up the best rate, whereby the burden imposed on the agent can be mitigated.

Also, if the best rate of the chained issue changes to such an extent as to become outside the designation range or

if a sell order for the chained issue is canceled, the chain order is deleted from the screens and the display of the chain order is suspended until the chain order again fulfills the display conditions thereafter. It is therefore possible to prevent a
5 buy order from being placed by mistake.

In the foregoing description, only a chain order in which one purchases somebody else's goods on condition of sale of his/her bonds is explained, but the present invention is also applicable to a chain order in which one sells his/her goods on
10 condition of purchase of somebody else's bonds.

Also, although the above description is given taking bonds as an example, goods to be transacted by the present invention are not limited to bonds alone and various other goods can of course be transacted.

15 Further, in the foregoing description, a buy order (or sell order) for chained goods is temporarily displayed on the board screens after a chain deal is conducted, but may be displayed from the first on the board screens while explicitly indicating that the goods are chained goods.

20 Referring now to FIG. 11, a second embodiment of the present invention will be described. In the figure, identical reference numerals are used to denote identical elements appearing in FIG. 5 and description of such elements is omitted.

In this embodiment, goods to be dealt in are "used
25 cars." Terminal units 20 and 24-1 to 24-3 each comprise a personal computer or the like. The terminal unit 20 is installed at broker's side while the terminal units 24-1 to 24-3 are

installed at agencies' side.

A network 21 is, for example, a public network such as ISDN (Integrated Service Digital Network), and interconnects the transaction assisting apparatus 1 with the terminal units
5 24-1 to 24-3, etc. to permit transmission of information from one to another.

A CAFIS (Credit And Finance Information System; trademark) 22 takes charge of a process for relaying credit data to a credit card company when payment is made by a credit card
10 or the like, and also provides services on behalf of a credit card company when the operation of the company's computer center is suspended.

A CMS (Cash Management Service) system 23 performs, among other processes, transfer of payment from one bank to
15 another when payment is made from a bank account.

The other elements are identical with those appearing in FIG. 5.

The operation of the embodiment will be now described.

FIG. 12 illustrates a replacement order which is
20 transacted in this embodiment. Replacement order is an order in which one requests, on condition of sale or purchase of certain goods, purchase or sale of other goods. In the example shown in FIG. 12, used cars are taken as an example of such goods and A denotes an orderer who placed a replacement order. B and C
25 are a purchaser and a seller, respectively, involved in the replacement order. Specifically, the replacement order placed by A is an order wherein A requests, on condition of sale of

his/her vehicle A or purchase of a vehicle B, purchase of the vehicle B or sale of his/her vehicle A. In the following, the scheme and process for appropriately transacting such a replacement order will be described.

5 FIG. 13 is a flowchart showing an example of a process executed when a sell/buy order for a used car is input in the embodiment shown in FIG. 11. Upon start of the process shown in the flowchart, the following steps are executed.

10 [S50] The CPU 1a accepts input of orders (ordinary orders and replacement orders) from the terminal units 20 and 24-1 to 24-3 via the I/Fs 1e and 1f.

15 FIG. 14 shows an example of an order (in the illustrated example, a replacement order) input from a terminal unit. In the figure, "ORDER NO." is a serial number uniquely assigned to each order, and "BODY NO." is a number specifying a vehicle model. "MODEL YEAR" indicates the year in which the vehicle was produced, and "COLOR" indicates the color of the vehicle body. "DISTANCE" indicates the distance traveled by the vehicle. "FIRST REGISTRATION DATE" indicates the date on which the vehicle was
20 purchased, and "LAST REGISTRATION DATE" indicates the date of the last vehicle inspection. "LAST REGISTRATION NO." indicates the vehicle number assigned to the vehicle at the time of the last registration, and "PRICE DESIRED" indicates a desired selling price of the vehicle. "DEAL" indicates the type of deal
25 and, in the illustrated example, "REPLACEMENT" is shown to the right of "DEAL", specifying that the order concerned is a replacement order. A summary of a replacement vehicle is shown

in the right-hand columns including the items "BODY NO." to
"PRICE". Specifically, the replacement vehicle should be as
follows: body number: "111845"; model year: "'89"; color:
"WHITE"; distance traveled: "80,000 km"; and price: "900,000"
5 yen.

"PRIORITY" indicates which of buy order and sell order
should be given priority. Namely, if buy order is to be given
priority, a procedure for selling the vehicle is taken after the
purchase of a desired vehicle. Conversely, if sell order is to
10 be given priority, a procedure for purchasing a desired vehicle
is taken after the sale of the vehicle.

"BANK NO." through "ACCOUNT HOLDER'S PHONE NO." show
information about a bank account from which payment is to be
transferred when the deal is conducted.

15 In the above example, payment is made from a bank
account, but may alternatively be made by a credit card. FIG.
15 shows an example of an order wherein payment is to be made
by a credit card. In this example, "CARD NO." through
"CARDHOLDER'S PHONE NO." show information necessary for the
20 payment by a credit card.

In the case where the order placed is an ordinary order
and not a replacement order, "ORDINARY" is entered as the type
of "DEAL" and the order form is formatted such that information
about a replacement vehicle ("PRIORITY" through "PRICE") is
25 omitted.

[S51] The CPU 1a stores the input order in a
predetermined area of the HDD 1d.

[S52] The CPU 1a determines whether or not the input order is a replacement order. If the input order is a replacement order, the flow proceeds to Step S54, and if not, the flow proceeds to Step S53.

5 [S53] The CPU 1a supplies the input order to the terminal units 20 and 24-1 to 24-3 via the I/Fs 1e and 1f to be displayed on their screen.

FIG. 16 shows an example of a screen displayed at the terminal units 20 and 24-1 to 24-3. In this example, the order
10 numbers "010001" through "010003" correspond to ordinary orders (not replacement orders). The contents of display appearing on the screen correspond to the individual items of orders input through the terminal units. In the figure, "SELL/BUY" indicates whether the order concerned is a sell order or a buy order.

15 [S54] The CPU 1a determines whether or not the replacement order specifies buy order to be given priority. If buy order is to be given priority, the flow proceeds to Step S55; if not, the flow proceeds to Step S56.

In the example shown in FIG. 14, for example, "BUY"
20 is shown to the right of "PRIORITY". It is therefore judged that buy order is to be given priority, and the flow proceeds to Step S55.

[S55] The CPU 1a extracts a buy order from the replacement order, and supplies the buy order to the terminal
25 units 20 and 24-1 to 24-3 to be displayed on their screen.

In FIG. 16, the buy order extracted from the replacement order shown in FIG. 14 is displayed at the lowermost

row. Namely, the items appearing in the right-hand column of the "DEAL" section in FIG. 14 are shown in the row. The underscore of "010004" as the "ORDER NO." indicates that the order concerned is a replacement order.

5 [S56] The CPU 1a extracts a sell order from the replacement order, and supplies the sell order to the terminal units 20 and 24-1 to 24-3 to be displayed on their screen.

 According to the process described above, orders input from the terminal units are stored in the HDD 1d shown in FIG. 10 11 and are also displayed at the individual terminal units as a screen as shown in FIG. 16, for example.

 Referring now to the flowchart of FIG. 17, a process executed by the transaction assisting apparatus 1 shown in FIG. 11 when a stored order is transacted will be described. The 15 process shown in the flowchart is executed at predetermined intervals as an interruption process, for example, to transact orders stored in the HDD 1d. Upon start of the process shown in the flowchart, the following steps are executed.

 [S60] The CPU 1a sets a variable i for counting the 20 number of times the process is executed, to an initial value of "1".

 [S61] The CPU 1a searches the orders stored in the HDD 1d to acquire an ith order.

 [S62] The CPU 1a determines whether or not the acquired 25 order is a replacement order. If the acquired order is a replacement order, the flow proceeds to Step S64, and if not, the flow proceeds to Step S63.

[S63] The CPU 1a executes an ordinary process with respect to an ordinary order.

The ordinary process means a process wherein a buy or sell order that fulfills the conditions of the acquired sell or buy order, if any, is reserved and the payment is adjusted.

[S64] The CPU 1a determines whether or not the replacement order specifies buy order to be given priority. If buy order is to be given priority, the flow proceeds to Step S65; if not, the flow proceeds to Step S66.

10 [S65] The CPU 1a executes a "BUY ORDER PREFERENCE PROCESS" wherein, if a sell order is placed for the buy order, such a sell order is reserved and then the sell order of the replacement order is conducted.

15 Details of the process will be described later with reference to FIG. 18.

[S66] The CPU 1a executes a "SELL ORDER PREFERENCE PROCESS" wherein, if a buy order is placed for the sell order, such a buy order is reserved and then the buy order of the replacement order is conducted.

20 Details of the process will be described later with reference to FIG. 19.

[S67] The CPU 1a determines whether or not the transaction of the replacement order has been completed. If the transaction has been completed, the flow proceeds to Step S68; 25 if not, the flow proceeds to Step S69.

[S68] The CPU 1a executes an "ADJUSTMENT PROCESS" wherein payments due arising from the transaction and taxes

accompanying the transaction are calculated and money is collected from or transferred to the bank accounts etc. of the parties involved in the transaction.

Details of the process will be described later with
5 reference to FIG. 21.

[S69] The CPU 1a increments the value of the variable i by "1".

[S70] If the value of the variable i is smaller than or equal to a total number N of orders then stored in the HDD
10 1d, the CPU 1a reverts to Step S61 to repeat the same process as described above; otherwise the process is ended.

Referring now to FIG. 18, the "BUY ORDER PREFERENCE PROCESS" shown in FIG. 17 will be described in detail. Upon start of the process shown in the flowchart, the following steps are
15 executed.

[S80] The CPU 1a determines whether or not the buy order has been transacted. If the buy order has already been transacted, the flow proceeds to Step S85; if not, the flow proceeds to Step S81.

20 Namely, Step S85 is executed if a sell order was placed for the buy order and the deal of the buy order has already been completed, or more specifically, if, in FIG. 12, only a contract for purchase of C's vehicle B by A has been concluded.

[S81] The CPU 1a searches the HDD 1d to determine
25 whether or not there is a sell order fulfilling the conditions of the buy order. If, as a result of the determination, there exists an applicable sell order, the flow proceeds to Step S82;

if not, the flow returns to the original process.

[S82] The CPU 1a reserves the applicable sell order.

For example, if there exists a sell order satisfying the conditions of the buy order with the order number "010004" shown in FIG. 16, such a sell order is reserved.

[S83] The CPU 1a stops displaying the corresponding buy order.

For example, in the example shown in FIG. 16, the buy order with the order number "010004" disappears from the screens of the individual terminal units.

[S84] The CPU 1a causes the individual terminal units to display the sell order of the replacement order.

In the example shown in FIG. 14, for example, information about the orderer's vehicle (vehicle with the body number "554845"), which is a target of the sell order, is supplied to the individual terminal units to be displayed thereat. FIG. 20 shows an example of a screen which is displayed at the individual terminal units as a result of execution of this step. In the illustrated example, a sell order with the order number "010004" is shown, specifying the vehicle with the body number "554845".

[S85] The CPU 1a determines whether or not a buy order satisfying the conditions of the sell order is stored in the HDD 1d. If such a buy order exists, the flow proceeds to Step S86; if not, the flow returns to the original process.

For example, if there exists a buy order fulfilling the conditions of the sell order with the order number "010004"

shown in FIG. 20, the flow proceeds to Step S86.

[S86] The CPU 1a reserves the buy order.

[S87] The CPU 1a causes the individual terminal units to stop displaying the sell order of the replacement order.

5 In the example shown in FIG. 20, for example, the sell order with the order number "010004" disappears from the screens.

[S88] The CPU 1a deletes the replacement order from the HDD 1d.

According to the above process, transactions can be
10 carried out in a manner such that the buy order of a replacement order is conducted preferentially over the sell order of same.

Referring now to FIG. 19, the "SELL ORDER PREFERENCE PROCESS" shown in FIG. 17 will be described in detail. Upon start of the process shown in the flowchart, the following steps are
15 executed.

[S90] The CPU 1a determines whether or not the sell order has been transacted. If the sell order has already been transacted, the flow proceeds to Step S95; if not, the flow proceeds to Step S91.

20 Namely, Step S95 is executed if a buy order was placed for the sell order and the deal of the sell order has already been completed, or more specifically, if, in FIG. 12, only a contract for purchase of A's vehicle A by B has been concluded.

[S91] The CPU 1a searches the HDD 1d to determine
25 whether or not there is a buy order fulfilling the conditions of the sell order. If, as a result of the determination, there exists an applicable buy order, the flow proceeds to Step S92;

if not, the flow returns to the original process.

[S92] The CPU 1a reserves the applicable buy order.

[S93] The CPU 1a stops displaying the corresponding sell order.

5 [S94] The CPU 1a causes the individual terminal units to display the buy order of the replacement order.

 [S95] The CPU 1a determines whether or not a sell order satisfying the conditions of the buy order is stored in the HDD 1d. If such a sell order exists, the flow proceeds to Step S96; 10 if not, the flow returns to the original process.

[S96] The CPU 1a reserves the sell order.

[S97] The CPU 1a causes the individual terminal units to stop displaying the buy order of the replacement order.

 [S98] The CPU 1a deletes the replacement order from 15 the HDD 1d.

According to the above process, transactions can be carried out in a manner such that the sell order of a replacement order is conducted preferentially over the buy order of same.

Referring now to FIG. 21, details of the "ADJUSTMENT 20 PROCESS" shown in FIG. 17 will be described. Upon start of the process shown in the flowchart, the steps stated below are executed. In the following description, the deals shown in FIG. 12 are taken as an example.

[S100] The CPU 1a calculates an overpayment r1 of 25 weight and automobile taxes with respect to the vehicle B which A purchased from C.

Namely, weight tax and automobile tax are collected

according to a system wherein the amount required until the next vehicle inspection must be paid in advance, and accordingly, if the title to the vehicle changes as a result of the transfer of the ownership from C to A, then it means that C has overpaid the taxes for an amount taxable after the date of transfer of the title. The amount thus borne in advance is therefore calculated as the tax overpayment r1.

[S101] The CPU 1a notifies the CMS system 23 of the sum of the price of the vehicle B and the tax overpayment r1, to have the sum transferred from A's account to C's account.

[S102] The CPU 1a causes the CMS system 23 to transfer the sum of the broker's commission plus consumption tax and an acquisition tax imposed on A for the acquisition of the vehicle from A's bank account to the broker's bank account.

The acquisition tax is calculated based on a price obtained by applying depreciation to new car's price.

[S103] The CPU 1a calculates an overpayment r2 of weight and automobile taxes with respect to the vehicle A which B purchased from A.

[S104] The CPU 1a notifies the CMS system 23 of the sum of the price of the vehicle A and the tax overpayment r2, to have the sum transferred from B's account to A's account.

[S105] The CPU 1a causes the CMS system to transfer the sum of the broker's commission plus consumption tax and an acquisition tax from B's account to the broker's account.

According to the above process, taxes and prices are automatically calculated and also are automatically transferred

from one bank account to another, whereby broker's procedures can be expedited and taxes accompanying transactions can be automatically calculated and collected.

As described above, the second embodiment of the present invention permits reliable and speedy transactions of replacement orders in cases where goods such as automobiles are traded, for example. Further, taxes accompanying transactions can be accurately calculated and collected from the right person who has to pay the taxes, so that the burden imposed on the broker can be mitigated.

In the above embodiment, used cars are taken as an example of goods to be dealt in, but the goods to be dealt in may be real estate or other goods, for example.

Further, instead of collecting the price and the taxes from a bank account, the payment of the price and taxes may be settled by a credit card, for example. Such payment can be made by use of the CAFIS 22 shown in FIG. 11.

Lastly, the processing function described above can be performed by a computer. In this case, the contents of the process for performing the functions as the transaction assisting apparatus may be described in a program recorded on a computer-readable recording medium. The program is executed by a computer, whereupon the above-described process is performed by the computer. The computer-readable recording medium includes a magnetic recording device, a semiconductor memory, etc.

To bring the program to market, the program may be

stored in portable recording media, such as CD-ROM (Compact Disk Read Only Memory) or floppy disk, to be distributed, or the program may be stored in a storage device of a computer connected to a network so that it can be transferred to other computers
5 through the network. The program may be stored in a hard disk unit or the like of a computer, and when it is to be executed by the computer, the program is loaded into the main memory and executed.

As described above, according to the present invention,
10 a transaction assisting apparatus for assisting transactions in goods comprises chain order input means permitting input of a chain order, the chain order being an order requesting, on condition of sale or purchase of first goods, purchase or sale of second goods different from the first goods, order detecting
15 means for detecting an order for purchase or sale of the first goods with respect to which the chain order is placed, and chain order processing means, responsive to the detection of a buy or sell order with respect to a predetermined chain order by the order detecting means, for performing a selling or purchasing
20 process with respect to the first goods and performing a purchasing or selling process with respect to the second goods. Accordingly, chain deals can be transacted reliably and speedily, and the burden imposed on the agent etc. can be mitigated.

The foregoing is considered as illustrative only of
25 the principles of the present invention. Further, since numerous modifications and changes will readily occur to those skilled in the art, it is not desired to limit the invention to

